Knowledge in Norms: A Sketch

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- Require: "If x executes U, then f ought to attain."
- Express: "x knows the plan U."
- Require: "For x to execute U, x must know U."

Constants: Users

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Let \mathbf{Q} be first-order logic built on these ingredients. To \mathbf{Q} , we add a deontic operator \bigcirc , obtaining \mathbf{QD}^* .

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Let Q be first-order logic built on these ingredients. To Q, we add a deontic operator \bigcirc , obtaining QD^{*}. The logic QD^{*} includes the constant domain assumption.

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But: functions are goal-directed.

We must represent the end of a user plan.

A goal is a state of affairs, a condition of the world \dots a formula of Q!

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We could also add preconditions to a user plan, as in dynamic logic.

pre :
$$Plan \longrightarrow \mathbf{Q}$$

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Assumes: every user can execute every plan.

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Note: this perspective on applications of plans is fundamentally coalgebraic!

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 $w \models [x, U]\varphi \quad \Leftrightarrow \quad \text{for all } w' \in \operatorname{app}_w(x, U),$ we have $w' \models \varphi$.

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After x applies U, the end of U ought to hold. Compare: $\bigcirc [x, U] \operatorname{end}(U)$

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- Epistemic operators can be added as the situation requires, of course.

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With this starting point, one can work to represent functional knowledge.

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Clearly, these tasks must go hand-in-hand.

Concrete steps:

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- Incorporate proper functions.
 - Represent "designer", "proper", etc.
 - Norms for proper function.
- Include epistemic operator for practical reasoning?